

# A Coq Library for Mechanised First-Order Logic

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The Coq Workshop'22, Haifa, Israel, August 12th

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# Contributors



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## Background

- Merge of several developments concerned with first-order logic
- Published at several venues (CPP, ITP, IJCAR, LFCS, FSCD, TYPES, JAR, JLC, LMCS)
- Design of a core framework general enough to accommodate all results
- Import of developments based on earlier versions of the framework
- Developed in a fork of the Coq library of undecidability proofs (Forster et al. (2020))

```
https://github.com/dominik-kirst/  
coq-library-undecidability/tree/fol-library/theories/FOL
```

# Framework

Emerged over several projects with ideas from various contributors:

- Deep embedding of syntax, deduction systems, and semantics
- Combination of well-known techniques, most notably de Bruijn indices
- Tool support for easy interaction by external users

Took most inspiration from O'Connor (2009); Ilik (2010); Herbelin and Lee (2014); Han and van Doorn (2020); Laurent (2021)

## Framework: Syntax

Terms and formulas are represented as inductive types  $\mathfrak{T}$  and  $\mathfrak{F}$  over a signature  $\Sigma = (\mathcal{F}_\Sigma, \mathcal{P}_\Sigma)$ :

$$t : \mathfrak{T} ::= x_n \mid f \vec{t} \quad (n : \mathbb{N}, f : \mathcal{F}_\Sigma, \vec{t} : \mathfrak{T}^{|\mathcal{F}|})$$

$$\varphi, \psi : \mathfrak{F} ::= \perp \mid P \vec{t} \mid \varphi \rightarrow \psi \mid \varphi \wedge \psi \mid \varphi \vee \psi \mid \forall \varphi \mid \exists \varphi \quad (P : \mathcal{P}_\Sigma, \vec{t} : \mathfrak{T}^{|\mathcal{P}|})$$

- Syntax modular in type classes for binary connectives and quantifiers
- Common instances  $(\rightarrow, \forall)$  and  $(\rightarrow, \wedge, \vee, \forall, \exists)$  provided
- Availability of  $\perp$  regulated via type class flag
- De Bruijn indices encode the number of quantifiers shadowing their relevant binder
- Capture-avoiding instantiation  $t[\sigma]$  and  $\varphi[\sigma]$  for parallel substitutions  $\sigma : \mathbb{N} \rightarrow \mathfrak{T}$

# Framework: Syntax (Coq)

```
Context {sig_funcs : funcs_signature}.
```

```
Inductive term : Type :=  
  | var : nat -> term  
  | func : forall (f : syms), vec term (ar_syms f) -> term.
```

```
Context {sig_preds : preds_signature}.
```

```
Inductive falsity_flag := falsity_off | falsity_on.  
Existing Class falsity_flag.
```

```
Class operators := {binop : Type ; quantop : Type}.  
Context {ops : operators}.
```

```
Inductive form : falsity_flag -> Type :=  
  | falsity : form falsity_on  
  | atom {b} : forall (P : preds), vec term (ar_preds P) -> form b  
  | bin {b} : binop -> form b -> form b -> form b  
  | quant {b} : quantop -> form b -> form b.
```

# Framework: Deduction Systems

Proof rules are represented as inductive predicates relating a context  $\Gamma$  to a formula  $\varphi$ :

$$\begin{array}{cccc} & & \dots & \\ \frac{\Gamma[\uparrow] \vdash \varphi}{\Gamma \vdash \forall \varphi} \text{ AI} & \frac{\Gamma \vdash \forall \varphi}{\Gamma \vdash \varphi[t]} \text{ AE} & \frac{\Gamma \vdash \varphi[t]}{\Gamma \vdash \exists \varphi} \text{ EI} & \frac{\Gamma \vdash \exists \varphi \quad \Gamma[\uparrow], \varphi \vdash \psi[\uparrow]}{\Gamma \vdash \psi} \text{ EE} \\ & & \dots & \end{array}$$

- Quantifier rules use shifted contexts  $\Gamma[\uparrow]$  so that  $x_0$  acts as canonical free variable
- Trivialises structural properties like substitutivity and weakening
- Availability of classical rules regulated via type class flag
- Similar representation of sequent calculi and other systems

# Framework: Deduction Systems (Coq)

```
Context {sig_funcs : funcs_signature}.
```

```
Context {sig_preds : preds_signature}.
```

```
Reserved Notation 'A ⊢ phi' (at level 61).
```

```
Inductive peirce := class | intu.
```

```
Existing Class peirce.
```

```
Inductive prv : forall (ff : falsity_flag) (p : peirce), list form -> form -> Prop :=  
  | II {ff} {p} A phi psi : phi::A ⊢ psi -> A ⊢ phi --> psi  
  | IE {ff} {p} A phi psi : A ⊢ phi --> psi -> A ⊢ phi -> A ⊢ psi  
  | AllI {ff} {p} A phi : map (subst_form ↑) A ⊢ phi -> A ⊢ ∀ phi  
  | AllE {ff} {p} A t phi : A ⊢ ∀ phi -> A ⊢ phi[t..]  
  | Exp {p} A phi : prv p A falsity -> prv p A phi  
  | Ctx {ff} {p} A phi : phi el A -> A ⊢ phi  
  | Pc {ff} A phi psi : prv class A (((phi --> psi) --> phi) --> phi)  
where 'A ⊢ phi' := (prv _ A phi).
```

## Framework: Semantics

Tarski models  $\mathcal{M}$  are represented as a domain type  $D$  and symbol interpretations:

$$f^{\mathcal{M}} : D^{|f|} \rightarrow D \qquad P^{\mathcal{M}} : D^{|P|} \rightarrow \mathbb{P}$$

- Interpretation of terms and formulas based on assignments  $\rho : \mathbb{N} \rightarrow D$
- Term evaluation  $\hat{\rho} t$  defined recursively, main rule  $\hat{\rho}(f \vec{t}) := f^{\mathcal{M}}(\hat{\rho} \vec{t})$
- Formula satisfaction  $\rho \models \varphi$  defined recursively, main rule  $\rho \models P \vec{t} := P^{\mathcal{M}}(\hat{\rho} \vec{t})$
- Induces the logical entailment relation  $\Gamma \models \varphi$

# Framework: Semantics (Coq)

```
Context {domain : Type}.
```

```
Class interp := B_I  
  { i_func : forall f : syms, vec domain (ar_syms f) -> domain ;  
    i_atom : forall P : preds, vec domain (ar_preds P) -> Prop ; }.
```

```
Definition env := nat -> domain.
```

```
Context {I : interp}.
```

```
Fixpoint eval (rho : env) (t : term) : domain := match t with  
  | var s => rho s  
  | func f v => i_func (Vector.map (eval rho) v) end.
```

```
Fixpoint sat {ff : falsity_flag} (rho : env) (phi : form) : Prop := match phi with  
  | atom P v => i_atom (Vector.map (eval rho) v)  
  | falsity => False  
  | bin Impl phi psi => sat rho phi -> sat rho psi  
  | quant All phi => forall d : domain, sat (d :: rho) phi end.
```

## Framework: Axiom Systems

Concrete axiom systems  $\mathcal{A}$  are modelled as predicates of formulas over a specific signature.

For the example of Peano arithmetic (PA), we instantiate to the arithmetical signature

$$(0, S \_, \_ + \_, \_ \times \_; \_ \equiv \_)$$

and collect the usual axioms, with the induction scheme represented as all instances of

$$\varphi[0] \rightarrow (\forall x. \varphi[x] \rightarrow \varphi[S x]) \rightarrow \forall x. \varphi[x].$$

- Include fragments of PA like Robinson's Q, also several variants of ZF set theory
- Equality  $\equiv$  seen as axiomatised symbol of the signature rather than a logical primitive
- Axiom systems  $\mathcal{A}$  induce relatives deductive and semantic theories  $\mathcal{A} \vdash \varphi$  and  $\mathcal{A} \models \varphi$

# Framework: Tool Support

Tools presented at last year's Coq Workshop (Hostert et al. (2021)):

- HOAS-input language
  - ▶ Concrete formulas can be written with Coq binders instead of de Bruijn indices
  - ▶ Eases interaction with the syntax
- Proof mode (inspired by Iris proof mode, Krebbers et al. (2017))
  - ▶ Tactic and notation layer hiding the proof rules
  - ▶ Eases interaction with the deduction systems
- Reification tactic (employing MetaCoq, Sozeau et al. (2020))
  - ▶ Extracts first-order formulas from Coq predicates
  - ▶ Eases interaction with the semantics

# Framework: Tool Support (Proof Mode)

```
205  frewrite (ax_add_zero y).
206  fapply ax_refl.
207  - fintros "x" "IH" "y".
208  frewrite (ax_add_rec (σ y) x).
209  frewrite ("IH" y).
210  frewrite (ax_add_rec y x). fapply ax_refl.
211 Qed.
212
213 Lemma add_comm :
214   FAI ⊢ << ∀' x y, x ⊕ y == y ⊕ x.
215 Proof.
216   fstart. fapply ((ax_induction (<< Free x, ∀' y, x ⊕ y == y ⊕ x))).
217   - fintros.
218     frewrite (ax_add_zero x).
219     frewrite (add_zero_r x).
220     fapply ax_refl.
221   - fintros "x" "IH" "y".
222     frewrite (add_succ_r y x).
223     frewrite <- ("IH" y).
224     frewrite (ax_add_rec y x).
225     fapply ax_refl.
226 Qed.
227
228 Lemma pa_eq_dec :
229   FAI ⊢ << ∀' x y, (x == y) ∨ ¬ (x == y).
230 Proof.
231   fstart.
232   fapply ((ax_induction (<< Free x, ∀' y, (x == y) ∨ ¬ (x == y)))).
233   - fapply
```

```
1 goal
p : peirce
x, y : term
_____ (1/1)
FAI
"IH" : ∀ x0, x'[†] ⊕ x0 == x0 ⊕ x'[†]
_____
[[σ x ⊕ y == y ⊕ σ x]]
```

Messages Errors Jobs

<https://github.com/dominik-kirst/coq-library-undecidability/blob/fol-library/theories/FOL/Proofmode/DemoPA.v>

# Framework: Tool Support (Reification Tactic)

```
87 Proof.  
88 elim a using PA_induction.  
89 - represent.  
90 - eapply ieq_trans. 1:apply (add_zero_l (iS b)).  
91   apply ieq_congr_succ, ieq_sym, add_zero_l.  
92 - intros d IH.  
93   eapply ieq_trans. 1:apply (add_succ_l d (iS b)).  
94   apply ieq_congr_succ. eapply ieq_trans.  
95     + apply IH.  
96     + apply ieq_sym, add_succ_l.  
97 Qed.  
98  
99 Lemma add_comm a b : a i⊕ b i= b i⊕ a.  
100 Proof.  
101 elim a using PA_induction.  
102 - represent.  
103 - eapply ieq_trans.  
104   + apply (add_zero_l b).  
105   + apply ieq_sym, (add_zero_r b).  
106 - intros a' IH.  
107   eapply ieq_trans. 2:eapply ieq_trans.  
108   + apply (add_succ_l a' b).  
109   + apply ieq_congr_succ, IH.  
110   + apply ieq_sym, add_succ_r.  
111 Qed.
```

```
1 goal  
D' : Type  
I : interp D'  
D_fulfills : forall (f : form) (rho : env D'),  
              PAeq f -> rho ⊨ f  
a, b : D'  
  
representableP 1  $\frac{(1/1)}{\llbracket \text{fun } a0 : D \Rightarrow a0 \text{ i} \oplus b \text{ i} = b \text{ i} \oplus a0 \rrbracket}$ 
```

Messages Errors Jobs

<https://github.com/dominik-kirst/coq-library-undecidability/blob/fol-library/theories/FOL/Reification/DemoPA.v>

## Framework: Evolution

Forster, Kirst, Smolka (2019) at CPP'19:

- Concrete signature, small logical fragment, named variables
- Among the initial projects constituting the undecidability library

Forster, Kirst, and Wehr (2021) at LFCS'20/JLC'21:

- Arbitrary signature, both logical fragments, de Bruijn encoding
- Use of Autosubst 2 (Stark et al. (2019)) for de Bruijn boilerplate

Kirst and Larchey-Wendling (2020) at IJCAR'20/LMCS'22:

- Parametric in logical fragment, merged into undecidability library
- Refrains from Autosubst 2 mostly due to dependency on function extensionality

Kirst and Hermes (2021) at ITP'21/JAR'22:

- Compromise of previous developments, merged into undecidability library
- Still no explicit code generation with Autosubst 2 but identical design

## Framework: Comparison

<b>Development</b>	<b>Signature</b>	<b>Binding</b>	<b>(AI)-Rule</b>	<b>Weakening</b>
O'Connor	arbitrary	named	side-condition	n.a.
Ilik	monadic	locally-nameless	co-finite	easy
Herbelin et al.	dyadic	locally-named	side-condition	needs renaming
Han and van Doorn	arbitrary	de Bruijn	shifting	easy
Laurent	full	anti-loc.-namel.	shifting	easy
Our framework	arbitrary	de Bruijn	shifting	easy

# Content

## Overview:

- Many metamathematical results: completeness, undecidability, incompleteness
- Many interdependencies, based on the Coq library of undecidability proofs
- Many possible projects/collaborations: syntactic cut-elimination, Hilbert systems, Löwenheim-Skolem theorems, resolution, tableaux, constructible hierarchy, ...

## Shared methods:

- Constructive meta-theory where possible
- Synthetic approach to computability results

## Content: Completeness

In which situations does  $\Gamma \vDash \varphi$  imply  $\Gamma \vdash \varphi$ ?

Based on the publication Forster et al. (2021):

- Constructively extremely subtle topic, extensive related literature
- Model-theoretic semantics (Tarski/Kripke) yield connections to MP and LEM
- Fully constructive proofs for algebraic and dialogical semantics

## Content: Undecidability

Which decision problems of first-order logic are undecidable?

Library includes all common undecidability results:

- Validity, provability, satisfiability (Forster et al. (2019))
- Finite satisfiability (Kirst and Larchey-Wendling (2020))
- Strongest versions regarding binary signatures (Hostert et al. (2022))
- Several variants of PA and ZF (Kirst and Hermes (2021))
- Post's theorem on the arithmetical hierarchy (Kirst et al. (2022))

## Content: Incompleteness

Which axiom systems  $\mathcal{A}$  satisfy  $\mathcal{A} \vdash \varphi$  or  $\mathcal{A} \vdash \neg\varphi$  for all  $\varphi$ ?

Library exploiting the connection to undecidability:

- Incompleteness of several variants of PA and ZF (Kirst and Hermes (2021))
- Essential incompleteness of Q (Peters and Kirst (2022))
- Tennenbaum's theorem on computable models of PA (Hermes and Kirst (2022))

## Current Status: Overview

- Completed core framework ✓
- Main completeness, undecidability, and incompleteness results imported ✓
- Essential incompleteness, Tennenbaum's theorem, and Post's theorem pending ✓
- Signature transformations and further computability results planned to be imported ✗
- Total: about 25k lines of code (8500 spec, 15500 proofs, 1000 comments), 110 files

# Current Status: Structure

 mark-koch Fix Proofmode MinZF demo	✓ 2722b67 4 days ago	 History
..		
 Arithmetics	Rename Deduction -> ND to prepare for Sequent	2 months ago
 Completeness2	Start using Asimpl in places	5 days ago
 Deduction	Finish Kripke Completeness, work on atom substitution	10 days ago
 Incompleteness	Tarski Constructions ported	2 months ago
 Proofmode	Fix Proofmode MinZF demo	4 days ago
 Reification	Tarski Constructions ported	2 months ago
 Semantics	Add validity facts, remove "not_strong" preservation	10 days ago
 Sets	Rename Deduction -> ND to prepare for Sequent	2 months ago
 Syntax	Start using Asimpl in places	5 days ago
 Undecidability	Start using Asimpl in places	5 days ago
 Utils	Port FOLP reduction, refactor PCP decidabilities	2 months ago
 FragmentSyntax.v	Start using Asimpl in places	5 days ago
 FullSyntax.v	Start using Asimpl in places	5 days ago

# Current Status: Pending Contributions

Filters  Labels 9 Milestones 0 [New pull request](#)

 2 Open  1 Closed Author Label Projects Milestones Reviews Assignee Sort

 **WIP: Add further incompleteness results**  
#4 opened 4 days ago by bn-peters • Draft  4 tasks

 **Add PrenexNormalForm and ArithmeticalHierarchy**  
#3 opened 5 days ago by SohnyBohny

Filters  Labels 9 Milestones 0 [New issue](#)

 1 Open  0 Closed Author Label Projects Milestones Assignee Sort

 **Proofmode bugs**  
#2 opened 8 days ago by bn-peters

# Current Status: Activity

History for [coq-library-undecidability](#) / [theories](#) / [FOL](#)

Commits on Jul 22, 2022

**Fix Proofmode MinZF demo**

 **mark-koch** committed 4 days ago ✓



2722b67



Commits on Jul 21, 2022

**Remove Require in section**

 **JoJoDeveloping** committed 5 days ago ✓

Verified



f55bae2



**Start using ASimpl in places**

 **JoJoDeveloping** committed 5 days ago ✓

Verified



840596b



**Merge branch 'fol-library' of github.com:dominik-kirst/coq-library-un...** ...

 **mark-koch** committed 5 days ago ✓



89d4850



**Fix ProofMode**

 **mark-koch** committed 5 days ago



6d9e47b



Commits on Jul 20, 2022

**Working and somewhat efficient ASimpl tactic**

 **JoJoDeveloping** committed 6 days ago ✓

Verified



2922b59



## Future Plans

- 1 Finish importing the remaining developments
- 2 Possible round of refactoring (proof mode performance, falsity flags)
- 3 Decide on a plan how to integrate with the undecidability library
- 4 Follow the release cycle of the undecidability library, itself following Coq
- 5 Possible timeline: opam package for Coq 8.16, add to Coq CI for Coq 8.17
- 6 At any time: help new users get started and contribute their developments!

Thanks for listening!

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